

# A Novel Approach for Data Encryption Standard Algorithm

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**Abstract**—Now a day's providing Security for data is complicated task we have so many security methods that are implemented and deployed but out of them few are using and serving the needs of society. And we can't say that any algorithm is perfect and avoids threats. The main goal of any design of any encryption algorithm must be security against unauthorized attacks. Within the last decade, data in both the private and public sectors are increased which requires Availability, Authentication, Confidentiality, Integrity. In this paper we are considering The DES algorithm that defines the mathematical steps that transform original text (plain text) into a cipher text (secret code) and also transform the cipher text back to the original text. Here introduces a new method to enhance the performance of the Data Encryption Standard (DES) algorithm. This is done by replacing the 8/32 S-Box instead of 6/4 S-Box. The output of each S-Box undergoes AND and XOR operation before going to the permutation P. In this paper we also proposed a new operation Addition modulo instead predefined XOR operation applied during the 16 round of the standard algorithm.

**Index Terms**— DES, Encryption, Decryption, asymmetric cryptography, symmetric cryptography.

## I. INTRODUCTION

**Cryptography:** is usually referred to as the study of secret. It is probably most important aspect of communication security and is becoming increasingly important as a basic building block for computer security[3]. Data that can be read and understood without any special measures is called plaintext. The method of representing plaintext in such a way as to hide its substance is called encryption. Encrypting plaintext results in unreadable form is called cipher-text. The process of reverting cipher-text to its original plaintext is called decryption. There are two techniques used for data encryption and decryption, which are:

**Asymmetric Cryptography** in this sender and recipient use different keys then it is known as asymmetrical or public key cryptography. The key used for encryption is called the public key and the key used for decryption is called the private key.

**Symmetric Cryptography** If sender and recipient use the same key then it is known as symmetrical or private key cryptography as shown in Fig 1. It is always suitable for long data streams. One of the symmetric block encryption algorithm is DES[9].

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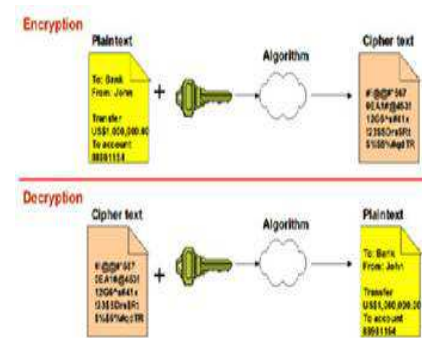


Fig. 1 Symmetric Key Encryption

Cryptography Goals: [2]

1. Confidentiality: Information in computer transmitted information is accessible only for reading by authorized parties.
2. Authentication: Origin of message is correctly identified with an assurance that identity is not false.
3. Integrity: Only authorized parties are able to modify transmitted or stored information.
4. Non-Repudiation: Requires that neither the sender, nor the receiver of message be able to deny the transmission.
5. Access Control- Requires access may be controlled by or for the target system.
6. Availability: Computer system assets are available to authorized parties when needed

## II. DATA ENCRYPTION STANDARD

The DES (Data Encryption Standard) algorithm is the most widely used encryption algorithm in the world [15]. For many years, and among many people, "secret code making" and DES have been synonymous. The DES was published by the United States' National Bureau of Standards in January 1977 as an algorithm to be used for unclassified data (information not concerned with national security)[1]. The Data Encryption Standard (DES), as specified in FIPS Publication 46-3, is a block cipher operating on 64-bit data blocks as shown in Fig 2

The encryption transformation depends on a 56-bit secret key and consists of sixteen Feistel iterations surrounded by two permutation layers: an initial bit permutation  $IP$  at the input, and its inverse  $IP^{-1}$  at the output. The structure of the cipher is depicted in Figure 2. The decryption process is the same as the encryption, except for the order of the round keys used in the Feistel iterations.[14] The 16-round Feistel network, which constitutes the cryptographic core of DES, splits the 64-bit data blocks into two 32-bit words, LBlock and RBlock (denoted by  $L_0$  and  $R_0$ ). In each iteration (or round), the second word  $R_i$  is fed to a function  $f$  and the result is added to the first word  $L_i$ . Then both words are swapped and the algorithm proceeds to the next iteration.

The function  $f$  of DES algorithm is key dependent and consists of 4 stages.

1. Expansion ( $E$ ): The 32-bit input word is first expanded to 48 bits by duplicating and reordering half of the bits.[7]
2. Key mixing: The expanded word is XORed with a round key constructed by selecting 48 bits from the 56-bit secret key, a different selection is used in each round.
3. Substitution. The 48-bit result is split into eight 6-bit words which are substituted in eight parallel 6x4-bit S-boxes. All eight S-boxes are different but have the same special structure.
4. Permutation ( $P$ ): The resulting 32 bits are reordered according to a fixed permutation before being sent to the output. The modified R Block is then XORed with LBlock and the resultant fed to the next RBlock register. The unmodified R Block is fed to the next L Block register. With another 56 bit derivative of the 64 bit key, the same process is repeated.

Pseudo Code: Data Encryption Standard:

INPUT: plaintext  $p_1 \dots p_{64}$ ; 64-bit key  $K=k_1 \dots k_{64}$  (includes 8 parity bits).  
 OUTPUT: 64-bit cipher text block  $C=c_1 \dots c_{64}$ .

1. (key schedule) Compute sixteen 48-bit round keys  $K_i$  from  $K$ .  
 Note: Where  $k=64$  bits out of which 8 parity bits are discarded outcome is 56 bits, after Left circular shift and PC2 which results 48 bit key.
2.  $(L_0, R_0) \leftarrow IP(p_1, p_2, \dots, p_{64})$  (Use IP Table to permute bits; split the result into left and right 32-bit halves  $L_0=p_{58}p_{50} \dots p_8, R_0=p_{57}p_{49} \dots p_7$ )
3. (16 rounds) for  $i$  from 1 to 16, compute  $L_i$  and  $R_i$  as follows:
  - 3.1.  $L_i=R_{i-1}$
  - 3.2.  $R_i=L_{i-1} \text{ XOR } f(R_{i-1}, K_i)$
 where  $f(R_{i-1}, K_i) = P(S(E(R_{i-1}) \text{ XOR } K_i))$ , computed as follows:
  - (a) Expand  $R_{i-1} = r_{32}r_{1r2} \dots r_{32}r_1$  from 32 to 48 bits,  $M \leftarrow E(R_{i-1})$ .
  - (b)  $M' \leftarrow M \text{ XOR } K_i$ . Represent  $M'$  as eight 6-bit character strings:  $M' = (B_1 \dots B_8)$
  - (c)  $M'' \leftarrow (S_1(B_1), S_2(B_2), \dots, S_8(B_8))$ . Here  $S_i(B_i)$  maps to the 4-bit entry in row  $r$  and column  $c$  of  $S_i$
  - (d)  $M''' \leftarrow P(M'')$ . (Use  $P$  per table to permute the 32 bits of  $M'' = m_1m_2 \dots m_{32}$ , yielding  $m_{16}m_7 \dots m_{25}$ )
4.  $b_1b_2 \dots b_{64} \leftarrow (R_{16}, L_{16})$ . (Exchange final blocks  $L_{16}, R_{16}$ .)
5.  $C \leftarrow IP^{-1}(b_1b_2 \dots b_{64})$ .
6. End.

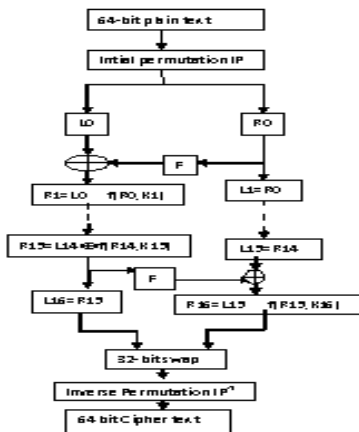


Fig. 2 DES Algorithm

III. PROPOSED DES ALGORITHM

In this paper we proposed a new improvement to the DES algorithm which makes the use of the new operation known as addition modulo (+).it takes two inputs and performs Addition and resulting output assume like  $x$ . later perform  $x \text{ mod } 2^w$  Where  $w$  is the number of bits that depends on given input.

Example:  $x$  and  $y$  are the Inputs

$X=1100\ 1000$

$Y=1000\ 1111$

$X^1$  is obtained by performing  $x+y$

$X^1 = 1\ 0101\ 0111$



Carry can be thrown off (or) perform modulo  $2^8$

$X^1$  is converted to decimal number

$X^1 = 343 \text{ mod } 2^8 = 87$

Binary equivalent of  $x^1$  is 0101 0111

To find original  $x$  value perform following operation

$X = x^1 + (-y)$

To obtain  $(-y) = 2^8 - y = 256 - 143 = 113$

Perform  $X^1 + (-y)$  which results

Original  $x$

0111 0001

0101 0111

1100 1000 ← original  $x$  value

The graphical representation of the proposed Des as shown in Fig.3

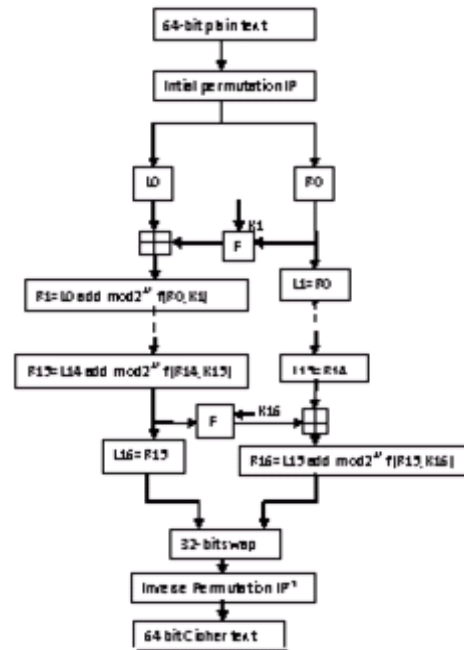


Fig. 3 Proposed DES Algorithm.

A. Algorithm of modified data encryption standard with addition modulo operation

INPUT: plaintext  $p_1 \dots p_{64}$ ; 64-bit key  $K=k_1 \dots k_{64}$  (includes 8 parity bits).

OUTPUT: 64-bit cipher text block  $C=c_1 \dots c_{64}$ .

1. (key schedule) Compute sixteen 48-bit round keys  $K_i$  from  $K$ .

Note: Where k=64 bits out of which 8 parity bits are discarded outcome is 56 bits, after Left circular shift and PC which results 48 bit key.

2.  $(L_0, R_0) \leftarrow IP(p_1, p_2, \dots, p_{64})$  (Use IP Table to permute bits; split the result into left and right 32-bit halves  $L_0=p_{58}p_{50} \dots p_8, R_0=p_{57}p_{49} \dots p_7$ )

3. (16 rounds) for i from 1 to 16, compute  $L_i$  and  $R_i$  as follows:

3.1.  $L_i=R_{i-1}$

3.2.  $R_i = L_{i-1}$  addition modulo  $2^{32}$   $f(R_{i-1}, K_i)$

where  $f(R_{i-1}, K_i) = P(S(E(R_{i-1}) XOR K_i))$ , computed as follows:

(a) Expand  $R_{i-1} = r_{32}r_{1r2} \dots r_{32} r_1$  from 32 to 48 bits,  $M \leftarrow E(R_{i-1})$ .

(b)  $M' \leftarrow M XOR K_i$ . Represent  $M'$  as eight 6-bit character strings:  $M' = (B_1 \dots B_8)$

(c)  $M'' \leftarrow F'$  where function  $F' = (((s_1 \wedge s_2) XOR s_3) \wedge s_4) XOR s_5 \wedge s_6$ . Here  $s_i(B_i)$  maps to the 8/32 S-Box that consist of 256 entries.

(d)  $M''' \leftarrow P(M'')$ . (Use P per table to permute the 32 bits of  $M'' = m_1m_2 \dots m_{32}$ , yielding  $m_1m_6m_7 \dots m_{25}$ )

4.  $b_1b_2 \dots b_{64} \leftarrow (R_{16}, L_{16})$ . (Exchange final blocks  $L_{16}, R_{16}$ .)

5.  $C \leftarrow IP^{-1}(b_1b_2 \dots b_{64})$ .

6. End.

During each round i, the right half of the block is expanded to 48 bits and XORed with a 48 bit internal key  $K_i$  derived from  $K$ . the result then passes through 6 s-boxes which are nonlinear substitutions results 32 output bits from 8 input bits.

The 32 bit result undergoes AND, XOR operations and is then permuted and performs addition modulo  $2^{32}$  with left half of the block. As shown in Fig.4 Finally the two halves of the block are swapped before going through the next round.

After completing 16 rounds the result will be undergone for  $IP^{-1}$  and finally we will get cipher text of 64 bits length. For the input given above the following TABLE.I shows 16 rounds and resulting output.

Input: 0123456789ABCDEF

Output: E46037E66BA9FEB8

TABLE I: Result of proposed DES 16 Rounds

Round	Sub Key( $K_i$ )	Left Bits( $L_i$ )	Right Bits( $R_i$ )
IP		CC00CCFF	F0AAF0AA
1	1B02EFFC7072	F0AAF0AA	CE620D18
2	79AEB9DBC9EF	CE620D18	22D1F6C2
3	55FC8A426F99	22D1F6C2	DE6B4E29
4	72ADD6DB351D	DE6B4E29	AB5BFB24
5	7CEC07EB53AE	AB5BFB24	E4977040
6	63A53E507B2F	E4977040	2BB88B30
7	EC84B7F618BC	2BB88B30	2531F060
8	F78A3AC13BFB	2531F060	D0B8BCB2
9	E0DBEBEDE781	D0B8BCB2	7436155C
10	B1F347BA464F	7436155C	56FFC0E2
11	215FD3DED386	56FFC0E2	D66AAA00
12	7571F59467E9	D66AAA00	587FC4EA
13	97C4D1FABA41	587FC4EA	A6FFD004
14	5F43B7F2E73A	A6FFD004	FAA3CA30
15	BF918D3D3F0A	FAA3CA30	E9FFF05C
16	EB3D8B0E17F5S	E9FFF05C	5BC44D34

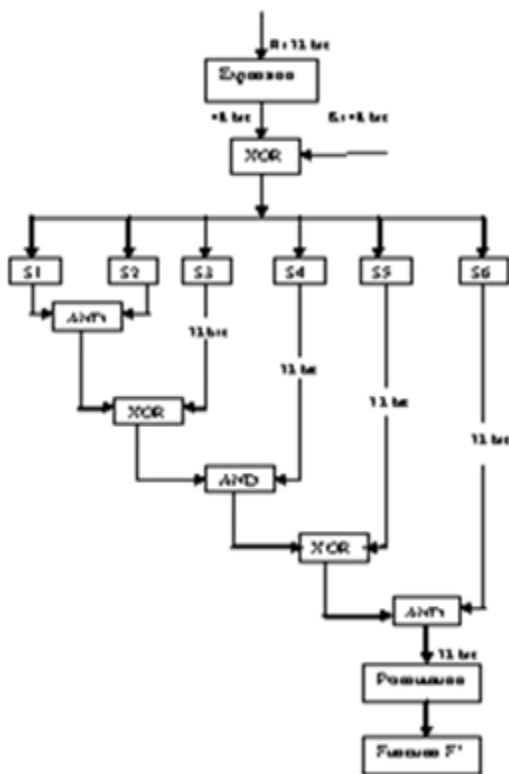


Fig. 4 Function F Design for proposed Algorithm

Here, using this proposed algorithm solve example. Our Input Message is 0123456789ABCDEF which is our plain text is converting into cipher text using this proposed algorithm. Here, the modified DES encrypts 64-bit blocks with a 56 bit key  $K$ . after an initial permutation of the bits, a plaintext block goes through 16 iterations (rounds) of a complex function and then passes through a final permutation that yields the cipher text block.

### B. Avalanche Effect

A change in one bit of the plain text or one bit of the key should produce a change in many bits of the cipher text this is referred to as the Avalanche Effect.

The following example shows the result when the first bit of the plain text is changed that results more changes in the resulting bits of the cipher text.

Plain text: 1123456789ABCDEF

↑  
Altered Bit

TABLE II: Result of avalanche effect for proposed DES 1stRound

Round	Input	Left Bits(Li)	Right Bits(Ri)	$\delta$
IP	0123456789A BCDEF11234 56789ABCD EF	CC00CCFF CC01CCFF	F0AAF0AA F1AAF0AA	2
1	CC00CCFF 0AAF0AACC 01CCFF1A AF0AA	F0AAF0AA F1AAF0AA	CE620D18 CE1CCE19	6

### IV. CONCLUSION

As we towards a society where automated information resources are increased and cryptography will continue to increase in importance as a security mechanism. DES is now considered to be insecure for some applications like banking system. There are also some analytical results which demonstrate theoretical weaknesses in the cipher. So it becomes very important to augment this algorithm by adding new levels of security to make it applicable. By adding modified S-Box design, modifies function implementation and replacing the old XOR by a new operation as proposed by this thesis to give more robustness to DES algorithm and make it stronger against any kind of intruding. This new algorithm gives avalanche effect than the original DES algorithm and also solves cryptanalysis attack

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